

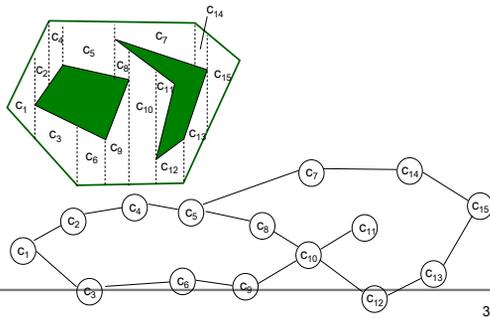
Path Quality

Spring 2013

Software Workshop:
High-Quality Motion Paths for Robots (and Other Creatures)

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Shortest paths among obstacles in the plane



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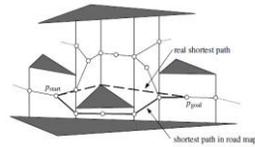
Today's lesson

- shortest paths
- high clearance paths
- other quality measures
- combined quality criteria and corridor maps
- path quality in sampling-based planners

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Shortest paths among obstacles in the plane [from de Berg et al, Ch. 15]

- first attempt: Dijkstra on the connectivity graph of the trapezoidal map



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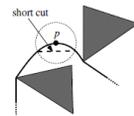
lesson

- does the graph on which we are searching for the best paths contain the best paths?

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properties of the shortest path

a polygonal line whose vertices are the start and goal configurations and vertices of the obstacles



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Computing a shortest path

Algorithm SHORTESTPATH(S, p_{start}, p_{goal})

Input. A set S of disjoint polygonal obstacles, and two points p_{start} and p_{goal} in the free space.

Output. The shortest collision-free path connecting p_{start} and p_{goal} .

1. $\mathcal{G}_{vis} \leftarrow \text{VISIBILITYGRAPH}(S \cup \{p_{start}, p_{goal}\})$
2. Assign each arc (v, w) in \mathcal{G}_{vis} a weight, which is the Euclidean length of the segment \overline{vw} .
3. Use Dijkstra's algorithm to compute a shortest path between p_{start} and p_{goal} in \mathcal{G}_{vis} .

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Computing the visibility graph

Algorithm VISIBILITYGRAPH(S)

Input. A set S of disjoint polygonal obstacles.

Output. The visibility graph $\mathcal{G}_{vis}(S)$.

1. Initialize a graph $\mathcal{G} = (V, E)$ where V is the set of all vertices of the polygons in S and $E = \emptyset$.
2. **for** all vertices $v \in V$
3. **do** $W \leftarrow \text{VISIBLEVERTICES}(v, S)$
4. For every vertex $w \in W$, add the arc (v, w) to E .
5. **return** \mathcal{G}

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Shortest paths in the plane, complexity

- the visibility-graph algorithm takes $O(n^2 \log n)$ time where n is the number of obstacle vertices
- there are output sensitive algorithms (in the size of the visibility graph)
- near-optimal $O(n \log n)$ algorithm by Hershberger and Suri
- the case of a simple polygon (whose complement is the obstacle) is much simpler

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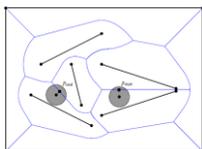
Shortest paths among polyhedra in 3-space

- the setting: point robot moving among polyhedra with a total of n vertices
- the problem is NP-hard [Canny-Reif]
 - algebraic complexity
 - combinatorial complexity

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High clearance paths

- Voronoi diagrams/the medial axis
- the Voronoi diagram of line segments, and the retraction method for a disc [O'Dunluing-Yap]



- video [Schirra/Rohnert]

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Other quality measures

- other L_p metrics, e.g., Manhattan (L_1)
- link number
- number of reverse movements
- low energy
- weighted regions
- many more

- multiple criterion optimal paths

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Combined quality criteria and corridor maps

- a path is called **Pareto optimal** if no other path has a better value for one criterion without having a worse value for another criterion
- multiple criterion optimization is often hard

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Clearance-length combination

- combined measure
- relaxed combination: the visibility Voronoi complex
- corridor maps

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Optimizing a combined measure

[Wein-van den Berg-H]

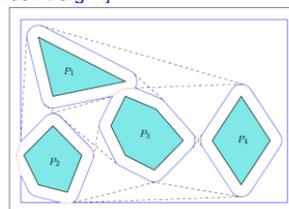
$$L^*(C) = \int_{\gamma} \left(\frac{w_{\max}}{w(t)} \right)^{d-1} dt$$

- examples:
 - the optimal path in the presence of a point obstacle is a logarithmic spiral
 - the optimal path in the presence of a segment obstacle is a circular arc
- approximation algorithms for the general polygonal case

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The visibility Voronoi diagram (VVD)

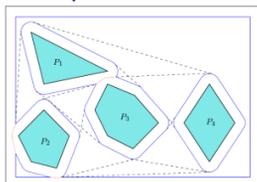
[Wein-van den Berg-H]



- finding the shortest path with a given clearance c , while still allowing to make significant shortcuts with lesser clearance on the Voronoi diagram

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The visibility Voronoi complex



- implicitly encodes the VVD for any clearance c
- interpolates between the visibility diagram ($c=0$) and the Voronoi diagram ($c=\infty$)
- $O(n^2 \log n)$ construction time

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Corridor maps

[Geraerts-Overmars]



- motivated by motion planning in games
- similar to VVD/VVC, augmenting the VD with clearance information
- instead of providing a single solution path, provides a **corridor** among static obstacles, where later one can easily maneuver among dynamic obstacles

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Path quality in sampling-based planners

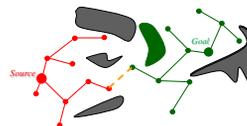
- the typical process: building a roadmap graph and running Dijkstra or similar
- recall our earlier test: **does the graph on which we are searching for the best paths contain the best paths?**
- path quality can be very low
- example: path length in BiRRT

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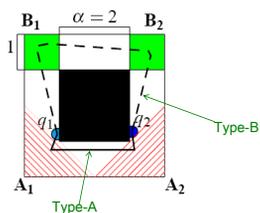
Bi-RRT reminder: Growing two-trees

[Kuffner and LaValle '00]

- maintain two trees rooted at **source** & **goal**
- construction step** – sample configurations and expand either tree as in RRT
- merging step** – connect configurations from both trees

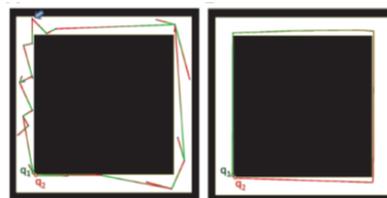


Example (I) – in OOPSMP



- 49.4% of paths are **over three times** worse than optimal (even after smoothing)
- much larger than the theoretical bound

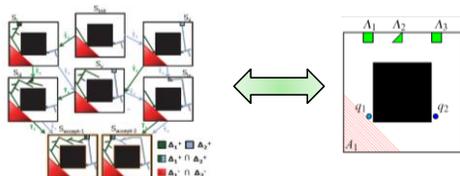
Example (II) – close-by start and goal configurations



- 5.9% of paths are **over 140 times worse** than optimal (even after smoothing)
- importance of **visibility blocking** – narrow passages not the only king (theoretical motivation for Visibility PRM, Laumond *et al.* '00)

How low can path quality get?

Sampling-Diagram Automata:
Analysis of path quality in tree planners
[Nechushtan-Raveh-Halperin, WAFR 2010]



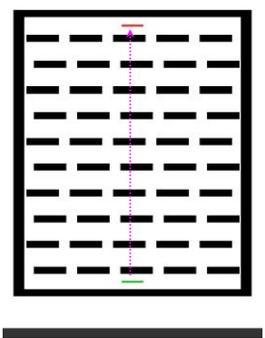
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Improving path quality in sampling-based motion planning, sample work

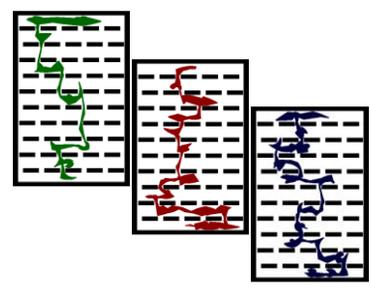
- Short-cutting heuristics ("path smoothing")
- Retraction towards medial axis [e.g., Wilmarth *et al.* '99, Geraerts and Overmars '07]
- Useful Cycles in PRM [Nieuwenhuisen and Overmars '04]
- Biasing tree growth by a cost-function [e.g., Urmson and Simmons '03, Ertlin and Bleuler '06, Jaillet *et al.* '08, Raveh *et al.* '09]
- RRT* - a modification of RRT [Karaman and Frazzoli '10] (for more variants, see paper)
 - the modified **RRT*** algorithm converges to an optimal path as running time reaches infinity
 - "Standard"-RRT **misses** the (precise) optimal path with probability one **Still, might be ϵ -good, or within same homotopy class as optimal path**

Improving quality by path hybridization

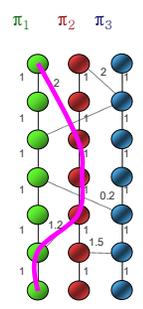
example: move the rod from the bottom to the top of a 2D grid (rotation + translation)



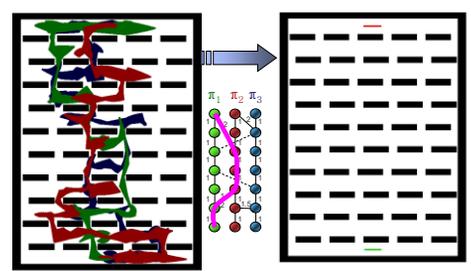
3 randomly generated motion paths



H-Graphs: Hybridizing multiple motion paths (= looking for shortcuts)



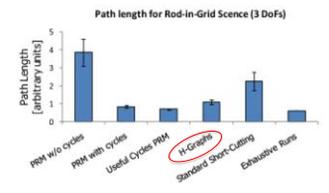
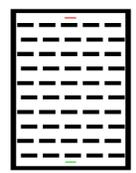
Hybridizing the paths



General quality criteria

Quality Measure	The Input Paths	H-Graph	Output Path
Clearance and length (emphasis on clearance)			
Clearance and length (emphasis on length)			
Path length			

Rod-in-Grid scene: 3 dof



Implemented in the **OOPSMP** package (Plaku, Moll and Kavraki), collision detection – **PQP** (Lin and Manocha)

Double-Wrench: 12 dof

Switching the two wrenches (rotation + translation x 2)

long runs of PRM
same time as total time of
HGraphs

Method	k-Inverse Clearance (arbitrary units)	Successful Runs %
PRM w/o cycles	~1.E+10	30%
PRM with cycles	~1.E+15	100%

H-Graphs become particularly useful for high-dimensional problems (at least in this example)

Scene adapted from Nieuwenhuisen et al., ICRA 04

Running-time bottleneck for hybridization:

Trying to connect nodes from different paths

in a naïve implementation:
 $O(n^2)$ potential edges need to be tested

Simple Heuristic – “Neighborhood H-Graphs”: compare only to nodes in local neighborhood – but can we do better?

we assume that paths are sorted by their start and end points

Edit-distance string matching

→ Linear alignment of motion paths

Comparing “This dog” and “That Dodge” with insertion / deletions / replacement:

THI - S DO - G -
THAT - DODGE

Alignment length is linear

Now testing only $O(n)$ edges along the alignment

dynamic-programming algorithm:

- insertion
- ↓ deletion
- ↘ replacement

		t	h	i	s	
h	0	1	1	1	1	4
a	1	1	1	1	1	3
s	2	1	1	1	1	3
	3	1	1	1	1	2

Comparison of running times

- hybridizing five motion paths in a 2-D maze:
 - from 3.52 seconds to 0.83 seconds on average (75% decrease), with comparable path quality

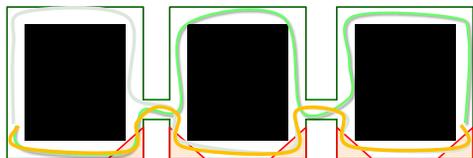
Method	Running Time (s)
H-Graphs	3.52
Edit Distance H-Graphs	0.83

IMPROVING THE QUALITY OF NON-HOLONOMIC MOTION BY HYBRIDIZING C-PRM PATHS

THOMAS BIEBER, FLORENCE COLLAS, DAN ZHANG, BRANKO RIKIĆ, DAN HALPERIN

applied to car-like motion with various quality criteria: length, smoothness, clearance, number of reverse vehicle motions

Why do Hgraphs work?



→ wrong decision can be taken at every step
→ can be solved by *path-hybridization*

References

- *Shortest Path and Networks*, J.S.B. Mitchell, Chapter 27 of the Handbook on DCG, Goodman-O'Rourke (eds)
- Visibility Graphs, Chapter 15 of the Computational Geometry book by de Berg et al
- more details, more experiments:
 - <http://acg.cs.tau.ac.il/projects>
 - *A Little More, A Lot Better: Improving Path Quality by a Path-Merging Algorithm* [Raveh-Enosh-Halperin] IEEE Trans. on Robotics, 2011

THE END