

# On the Number of Regular Vertices of the Union of Jordan Regions\*

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## Abstract

Let  $\mathcal{C}$  be a collection of  $n$  Jordan regions in the plane in general position, such that each pair of their boundaries intersect in at most  $s$  points, where  $s$  is a constant. Let  $U$  denote the union of  $\mathcal{C}$ . If the boundaries of two sets in  $\mathcal{C}$  cross exactly twice, then their intersection points are called *regular vertices* of the arrangement  $\mathcal{A}(\mathcal{C})$ . Let  $R(\mathcal{C})$  denote the set of regular vertices on the boundary of the union of  $\mathcal{C}$ . We present several bounds on  $|R(\mathcal{C})|$ , determined by the type of the sets of  $\mathcal{C}$ . (i) If each set of  $\mathcal{C}$  is convex, then  $|R(\mathcal{C})| = O(n^{1.5+\epsilon})$  for any  $\epsilon > 0$ .<sup>1</sup> (ii) If  $\mathcal{C}$  consists of two collections  $\mathcal{C}_1$  and  $\mathcal{C}_2$  where  $\mathcal{C}_1$  be a collection of  $n$  convex *pseudo-disks* in the plane (closed Jordan regions with the property that the boundaries of each pair of them intersect at most twice), and  $\mathcal{C}_2$  is a collection of convex polygons with a total of  $n$  sides, then  $|R(\mathcal{C})| = O(n^{4/3})$ , and this bound is tight in the worst case. (iii) If no further assumptions are made on the sets of  $\mathcal{C}$ , then we show that there is an integer  $t$  that depends only on  $s$  such that  $|R(\mathcal{C})| = O(n^{2-1/t})$ .

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<sup>1</sup>Throughout this paper,  $\epsilon$  stands for an arbitrarily small positive constant; the constants of proportionality in bounds that involve  $\epsilon$  also depend on  $\epsilon$  and generally, tend to  $\infty$  as  $\epsilon$  approaches zero.

## 1 Introduction

Let  $\mathcal{C}$  be a collection of  $n$  Jordan regions (the interiors of closed Jordan curves) in the plane, with the property that each pair of boundaries intersect in at most some constant number,  $s$ , of points. We assume that the sets of  $\mathcal{C}$  are in general position, that is no point is incident to more than two boundaries, and at each intersection point that the boundaries meet they cross. Let  $U$  denote the union of  $\mathcal{C}$ . We consider the arrangement  $\mathcal{A}(\mathcal{C})$ , formed by the boundaries of the sets in  $\mathcal{C}$ , and define a vertex of  $\mathcal{A}(\mathcal{C})$  (an intersection point between two boundaries) to be *regular* if the two boundaries cross exactly twice; all other vertices are called *irregular*. The goal is to obtain sharp bounds on the maximal number of regular vertices that appear on  $\partial U$ . We denote, as above, the set of regular vertices on  $\partial U$  as  $R(\mathcal{C})$ .

The interest in this problem goes back to the work of Kedem et al. [11], where it was shown that if all vertices of  $\mathcal{A}(\mathcal{C})$  are regular (such a collection  $\mathcal{C}$  is called a collection of *pseudo-disks*), then the number of (regular) vertices of  $\partial U$  is at most  $6n - 12$ , for  $n \geq 3$ , and this bound is tight in the worst case. Recently, Pach and Sharir [15] have shown that if  $\mathcal{C}$  is an arbitrary finite collection of convex sets (so that pairs of their boundaries can intersect in an arbitrary number of points) then one has  $|R(\mathcal{C})| \leq 2|I(\mathcal{C})| + 6n - 12$ , where  $I(\mathcal{C})$  is the set of irregular vertices on  $\partial U$ . This result was instrumental in a recent paper by Efrat and Sharir [6], showing that the complexity of the union of  $n$  planar “fat” convex sets, each pair of whose boundaries intersect in at most some constant number of points, is nearly linear in  $n$ . However, since  $I(\mathcal{C})$  can be  $\Omega(n^2)$ , for general collection  $\mathcal{C}$ , even when no pair of boundaries cross at more than four points the bound of [15] only yields the trivial  $O(n^2)$  upper bound on  $|R(\mathcal{C})|$ . An example of such a construction, imagine a collection of  $2n$  very narrow axis-parallel rectangles, from which  $n$  are vertical,  $n$  are horizontal, and each vertical one intersects each horizontal one at a point which is on the boundary of their union.

Pach and Sharir [15] also give a construction where the number of regular vertices on the boundary of the union of  $n$  rectangles and  $n$  congruent disks is  $\Omega(n^{4/3})$ . This is the best known lower bound for the general problem stated above, with a constant number of intersections between any pair of boundaries. Without this last constraint, it is easy to obtain examples with  $\Omega(n^2)$  regular vertices on the boundary of the union; see, e.g., [12, 15].

In Section 2 we show that the  $\Omega(n^{4/3})$  lower bound is tight for the special class (ii) of collections  $\mathcal{C}$  as in the abstract. Specifically, we show:

**Theorem 1.1** *The maximum possible number of regular vertices on the boundary of the union of a family of  $n$  convex pseudo-disks and of convex polygons with a total of  $n$  edges is  $\Theta(n^{4/3})$ .*

For the other two main results, we first present a technique for transforming the family of regions so that every regular vertex in  $R(\mathcal{C})$  becomes essentially a point of tangency between the two regions, and so that the number of intersections between any pair of boundaries does not increase (see Lemma 3.1).

Next, in Section 4, we consider the case of general convex regions, and show:

**Theorem 1.2** *The maximum possible number of regular vertices on the boundary of the union of a family of  $n$  convex Jordan regions in the plane where each pair of boundaries intersect in at most a constant number of points, is  $O(n^{1.5+\epsilon})$  for any  $\epsilon > 0$ .*

Finally, in Section 5, we study the case of general Jordan regions, and show:

**Theorem 1.3** *The maximum possible number of regular vertices on the boundary of the union of a family of  $n$  Jordan regions in the plane, where each pair of boundaries intersect in at most a constant number  $s$  of points, is  $O(n^{2-1/t})$ , where  $t$  is an integer constant that depends on  $s$ .*

In other words, we show that in a fairly general setting, the number of regular vertices to  $\partial U$  is subquadratic. This is likely to have implications on the analysis of the complexity of union of geometric objects in two and three dimensions (this is indicated by our current work in progress) and to shed more light on the combinatorial and topological structures of arrangements of Jordan regions in the plane. Our experience has been that improved combinatorial bounds on the complexity of the union of geometric objects often entails efficient algorithms for computing such boundaries. This in turn could have useful implications in applications like robot motion planning [9], solid modeling and others.

**Remark:** Note that using a straightforward perturbation scheme, it is clear that the maximum number of regular vertices on the union is achieved when the given regions are in general position, so this assumption involves no loss of generality.

## 2 Pseudo-disks and Polygons: Proof of Theorem 1.1

Let  $\mathcal{C} = \mathcal{C}_1 \cup \mathcal{C}_2$  where  $\mathcal{C}_1$  is a collection of  $n$  convex *pseudo-disks* in the plane (closed Jordan regions with the property that the boundaries of each pair of them intersect at most twice), and  $\mathcal{C}_2$  is a collection of convex polygons with a total of  $n$  sides. We form the union  $U_1$  of  $\mathcal{C}_1$ . By [11], the boundary of  $U_1$  consists of at most  $6n$  convex arcs, each lying on the boundary of a single pseudo-disk. Without loss of generality, we will treat only arcs that lie on the top boundary of their respective pseudo-disks. Let  $\Gamma$  denote the set of these arcs. Clearly, the relative interiors of the arcs in  $\Gamma$  are pairwise disjoint. Let  $m \leq 6n$  denote the size of  $\Gamma$ .

Let  $U_2$  denote the union of  $\mathcal{C}_2$ , and put  $U = U_1 \cup U_2$ . Clearly, the number of regular vertices on the boundary of  $U$  that are incident to two boundaries of sets in  $\mathcal{C}_1$  is at most  $6n$ .

It is also easy to show that the number of regular vertices that lie on the boundary of  $U$  and are incident to two boundaries of polygons in  $\mathcal{C}_2$  is  $O(n)$ . Indeed, let  $v$  be such a vertex, and let  $c, c'$  be the two polygons in  $\mathcal{C}_2$  whose boundaries contain  $v$ . Let  $K = c \cap c'$ . Recall that  $K$  must be a convex polygon whose boundary contains only two vertices (one of which is  $v$ ) where  $\partial c$  and  $\partial c'$  meet. Hence it must also contain at least one vertex  $w$  of  $c$  or of  $c'$ , such that  $vw$  is an edge of  $K$ . We charge  $v$  to such a vertex  $w$ , and note that any vertex can be charged at most twice in this manner. Hence the total number of vertices of  $\partial U$  of this type is  $O(n)$ . It thus remains to bound the number of *mixed* regular vertices of  $\partial U$ , namely, those that are incident to the boundaries of a pseudo-disk in  $\mathcal{C}_1$  and of a polygon in  $\mathcal{C}_2$ .

Let  $\gamma \in \Gamma$ , let  $d \in \mathcal{C}_1$  be the pseudo-disk of  $\mathcal{C}_1$  whose boundary contains  $\gamma$ , and consider its interaction with a polygon  $c$  of  $\mathcal{C}_2$  that has a regular vertex on  $\gamma$ . Either  $c$  contains the extreme left or right point of  $d$ , or  $c \cap \partial d$  lies completely in the top portion of  $\partial d$  (bottom portions are dealt with by a symmetric argument). Polygons  $c \in \mathcal{C}_2$  that contain the leftmost (rightmost) point of  $d$  contribute at most one mixed vertex of  $\partial U$  on  $\gamma$ , for a total of at most  $2n$  vertices. Similarly we can eliminate polygons  $c$  which contain one of the endpoints of  $\gamma$ , as those produce at most  $2m \leq 12n$  mixed vertices on  $\partial U$ . From this point on we restrict our

attention to polygons  $c$  such that  $c \cap \partial d \subset \gamma$ . Let  $c_1, \dots, c_k$  be the sequence of polygons in  $\mathcal{C}_2$  whose boundaries form with  $\gamma$  regular vertices on  $\partial U$ , arranged in the left-to-right order of these vertices along  $\gamma$  — observe that, if both intersections of  $\partial c_i$  with  $\gamma$  appear on  $\partial U$ , no other regular vertex on  $\gamma \cap \partial U$  can appear between them, so the above order is well defined. Let  $v, w$  be the two points of intersection of  $\gamma$  with  $\partial c_1$ , say. If  $v, w$  lie on different edges of  $c_1$  then the portion of  $c_1$  inside  $d$  must contain a vertex of  $c_1$ , and we can then charge  $v$  and  $w$  to such a vertex, in the same manner as in the preceding paragraph. It follows that the number of such vertices  $v, w$  is  $O(n)$ , and we can thus ignore such cases in our analysis. We can thus assume that  $\gamma$  crosses each  $\partial c_i$  at two points that lie on the same edge,  $e_i$ , of  $c_i$ . Let  $CH(e_i \cap \gamma)$  denote the convex hull of  $e_i \cap \gamma$ , namely the segment of  $e_i$  bounded by its two intersection points with  $\gamma$ . Clearly, the segments  $CH(e_1 \cap \gamma), \dots, CH(e_k \cap \gamma)$  are in convex position: each of them is fully contained in the lower envelope  $E$  of the lines containing these edges. We refer to these remaining regular vertices as *edge-touching* vertices.

We adapt the analysis technique of [7, 8]. First, we derive a weaker upper bound on the number of mixed edge-touching regular vertices on  $\partial U$ . To this end, we construct a bipartite graph  $H$  whose nodes are the arcs in  $\Gamma$  and the polygon edges, and each of its edges connects an arc  $\gamma$  to a polygon edge  $e$  if  $\gamma$  crosses  $e$  twice, at two regular vertices, at least one of which lies on  $\partial U$ . We claim that  $H$  does not contain a  $K_{2,3}$  as a subgraph (composed of two arcs in  $\Gamma$  and of three polygon edges). Indeed, suppose that  $H$  did contain such a subgraph, consisting of two arcs  $\gamma_1, \gamma_2 \in \Gamma$ , and of three polygon edges  $e_1, e_2, e_3$ . As noted, the segments  $CH(e_i \cap \gamma_1)$  are fully contained in the lower envelope of the three respective lines that contain these edges. Suppose, with no loss of generality, that  $e_1, e_2, e_3$  appear in this left-to-right order along the envelope; see Figure 1. The lower envelope is partitioned by  $\gamma_1$  into seven sections: the three segments  $CH(e_i \cap \gamma_1)$ ,  $i = 1, 2, 3$ , and four maximal complement sections of the envelope where the edge-touching vertices of  $\gamma_2$  and the edges  $e_i$  may appear. Since  $\gamma_1$  and  $\gamma_2$  are openly disjoint, the edge-touching vertices induced by  $\gamma_2$  can only appear along a single section of the latter four complement sections of the envelope. However, none of these sections contain portions of all three edges. Hence  $\gamma_2$  cannot connect to all of  $e_1, e_2$ , and  $e_3$  in the graph.

Since the graph  $H$  is bipartite, with at most  $m$  and  $n$  nodes, respectively, in each class, and since it contains no  $K_{2,3}$  as a subgraph, it follows from standard extremal graph-theoretic arguments (see [14]) that the number of its edges, and hence the number of mixed edge-touching regular vertices on  $\partial U$  is  $O(mn^{1/2} + n)$ .

We next choose an integer parameter  $r$ , to be fixed later, and construct a  $(1/r)$ -cutting on the edges of the polygons in  $\mathcal{C}_2$  (see [10] for details). This yields a tiling of the plane by  $O(r^2)$  pairwise openly disjoint trapezoids, each crossed by at most  $n/r$  edges. For each trapezoid  $\tau$ , consider the set  $\Gamma_\tau$  of all the arcs in  $\Gamma$  that cross  $\tau$ , clipped to within  $\tau$ ; some of these arcs may intersect  $\tau$  in two connected portions, and we regard each such portion as a separate arc.

We classify the arcs in  $\Gamma_\tau$  according to the pairs of sides of  $\tau$  that they cross. One class consists of *short* arcs that have at least one endpoint inside  $\tau$ ; any other, *long* arc meets  $\partial\tau$  exactly twice. One class of long arcs is referred to as the class of *bottom* arcs; these are the arcs that have both endpoints on the bottom side of  $\tau$ .

We claim that, for any class  $\Gamma'$  of arcs, other than those of the short arcs or of the bottom arcs, the number of mixed edge-touching regular vertices on  $\partial U$  that are formed within  $\tau$  by arcs of  $\Gamma'$  is  $O(n/r)$ . The proof is similar to that in [7]. It is based on the observation that none of the at most  $n/r$  polygon edges that cross  $\tau$  can form mixed edge-touching regular vertices

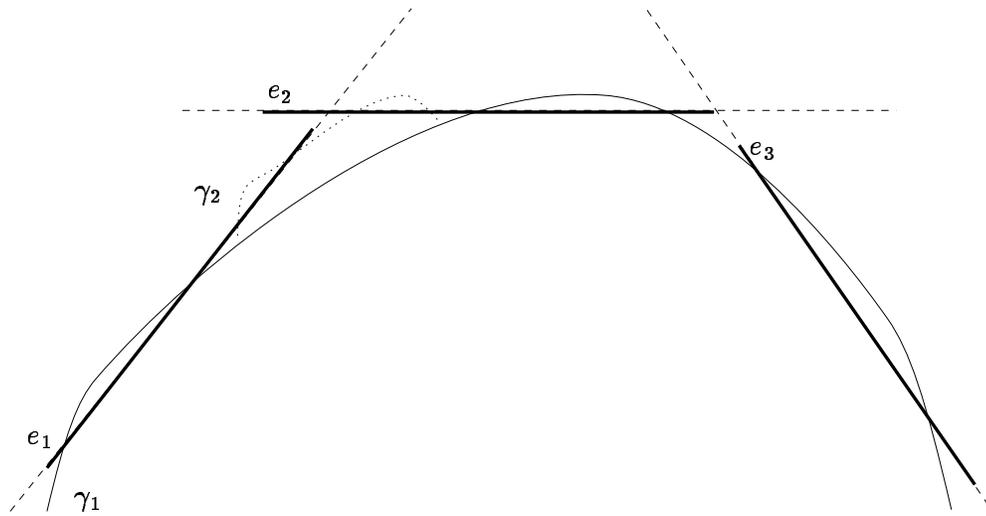


Figure 1: The graph  $H$  does not contain any  $K_{2,3}$  consisting of two arcs  $\gamma_1, \gamma_2$ , and three polygon edges  $e_1, e_2, e_3$

on  $\partial U$  with more than one arc of  $\Gamma'$ . Indeed, if this could happen for one such edge  $e$  and two such arcs  $\gamma, \gamma'$ , then  $CH(e \cap \gamma)$  and  $CH(e \cap \gamma')$  must be disjoint, with, say,  $e \cap \gamma$  lying to the left of  $e \cap \gamma'$ . (Indeed, the convexity of the pseudo-disks rules out the case that these two intervals are nested within each other; on the other hand, if these intervals overlap without being nested, then  $\gamma$  and  $\gamma'$  would have to cross, which is impossible.) If we trace  $\gamma$  from the right endpoint of  $e \cap \gamma$  to the right, and trace  $\gamma'$  from the left endpoint of  $CH(e \cap \gamma')$  to the left, then these curves must cross within  $\tau$ . Indeed, for this not to happen, either one of these curves has to end inside  $\tau$  (so  $\Gamma'$  is the class of short arcs), or both extensions must cross the bottom side of  $\tau$  (so  $\Gamma'$  is the class of bottom arcs). Since  $\Gamma'$  is neither of these classes, the claim follows; see [7, 8] for a similar argument. We have thus shown that the number of mixed edge-touching regular vertices on  $\partial U$  that are formed within  $\tau$  by arcs that are neither short nor bottom is  $O(n/r)$ , for an overall bound of  $O(r^2 \cdot (n/r)) = O(nr)$ , over all trapezoids  $\tau$ .

We next claim that the total number of bottom arcs, over all trapezoids  $\tau$ , is  $O(r^2 + m)$ . We show this using a graph-planarity argument, similar to that used in [8]. We construct a plane embedding of a planar graph  $G$  as follows. The nodes of  $G$  are the bottom edges of the trapezoids of our cutting. The edges of  $G$  are defined and drawn as follows.

Let  $\gamma'$  be a bottom subarc of some arc  $\gamma \in \Gamma$ , formed within some trapezoid  $\tau$ . Let  $e$  denote the bottom edge of  $\tau$ , and let  $u$  and  $v$  be the two intersection points of  $\gamma$  with  $e$ , where  $u$  lies to the left of  $v$  (these are the endpoints of  $\gamma'$ ). Now follow  $\gamma$  to the right from  $v$  until another bottom edge  $e'$  of some trapezoid  $\tau'$  is encountered; denote this subarc of  $\gamma$  by  $\gamma^R$ . Note that  $\gamma$  can hit  $e'$  either from above or from below. (If we do not meet any bottom edge, we can charge  $\gamma'$  uniquely to the right endpoint of  $\gamma$ , so the overall number of these bottom arcs is at most  $m$ .) We distinguish between two cases:

If the portion of  $\gamma$  to the left of  $u$  does not intersect  $e'$ , then we connect  $e$  and  $e'$  in the graph  $G$  along  $\gamma^R$ . If, on the other hand, the portion of  $\gamma$  to the left of  $u$  does intersect  $e'$ , then we claim that both  $u$  and  $v$  must lie on the top edge of  $\tau'$ . Indeed, by construction,  $\gamma^R$  does not

meet any nonvertical edge of the cutting. Moreover,  $\gamma^R$  must hit  $e'$  from above, for otherwise the convexity of  $\gamma$  implies that the entire portion of  $\gamma$  to the left of the hitting point lies below the line supporting  $e'$ , so it cannot meet  $e'$  again. Now if we follow  $\gamma^R$  from  $e'$  to the left, the first vertical edge of the cutting that we meet must be the left edge of  $\tau'$ , but then  $\gamma$  could not have met  $e'$  again further to the left. Hence  $\gamma^R$  is fully contained in  $\tau'$ , so  $v$  lies on the top side of  $\tau'$ , and a similar, slightly modified, argument implies the same for  $u$ . In fact, the same argument also implies that the portion of  $\gamma$  between  $u$  and  $e'$  is also fully contained in  $\tau'$ . Let  $u'$  and  $v'$  be the intersection points of  $\gamma$  with  $e'$ , where  $u'$  lies to the left of  $v'$ . We now apply the same analysis to the portions of  $\gamma$  extending to the left of  $u'$  and to the right of  $v'$ , respectively. We iterate this process until those portions do not end up on the same bottom side. When this is the case, we add the last right portion of  $\gamma$  as an edge of the graph  $G$ . (Again, if during this process we reach the right endpoint of  $\gamma$ , we charge  $\gamma'$  uniquely to that endpoint; the overall number of these bottom subarcs is at most  $m$ .) See Figure 2. Note that the entire portion of  $\gamma$  between (the final) points  $u'$  and  $v'$  contains only one bottom subarc, namely  $\gamma'$ . Hence the number of edges of  $G$  is equal to the number of bottom subarcs excluding those that we have already charged to the right endpoint of their containing arcs.

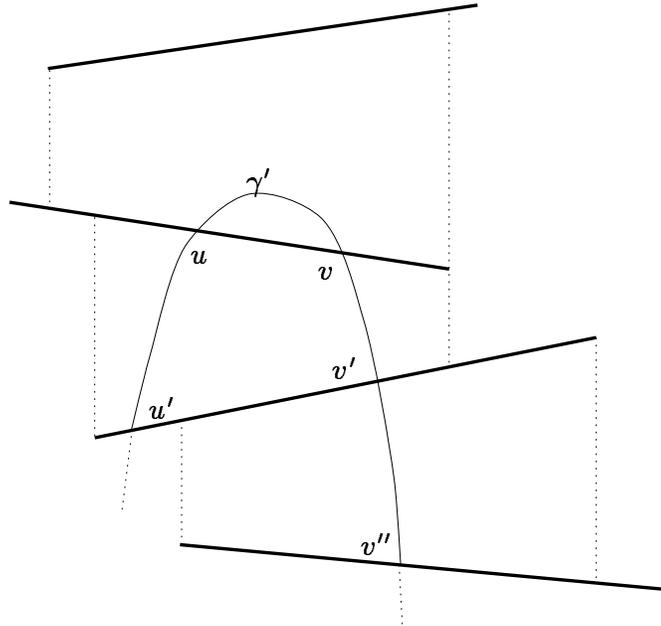


Figure 2: The portion of  $\gamma$  between  $v'$  and  $v''$  is added as an edge to the graph  $G$

Since the relative interiors of the arcs of  $\Gamma$  are pairwise disjoint,  $G$  is indeed a plane embedding of a planar graph. Let  $f$  be a face of  $G$  of degree 2. By construction,  $f$  is bounded by two subarcs  $\delta, \delta'$  of two respective distinct arcs  $\gamma, \gamma'$  of  $\Gamma$ , and by portions of two bottom edges  $e, e'$  of two respective trapezoids  $\tau, \tau'$ . If the interior of  $f$ , as a planar region, contains an endpoint of  $e$  or of  $e'$ , then we can charge  $f$  to this endpoint (see Figure 3 for an illustration of this case). Since there are  $O(r^2)$  such endpoints and each of them is charged at most once, the number of such faces  $f$  is  $O(r^2)$ . Otherwise, both  $\delta, \delta'$  must leave  $e$  from the same (top or bottom) side of it, and similarly for  $e'$ . (We say that  $\delta$  leaves the top (bottom) side of  $e$ , if in a small neighborhood of the intersection point of  $\delta$  and  $e$ ,  $\delta$  is on the top (bottom) side of  $e$ .) Since, by construction, each of  $\delta, \delta'$  must leave at least one of  $e, e'$  from its bottom side, there are only two possible cases:

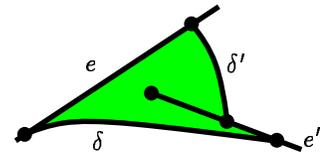


Figure 3: A face of  $G$  of degree 2 containing an endpoint of a bottom side

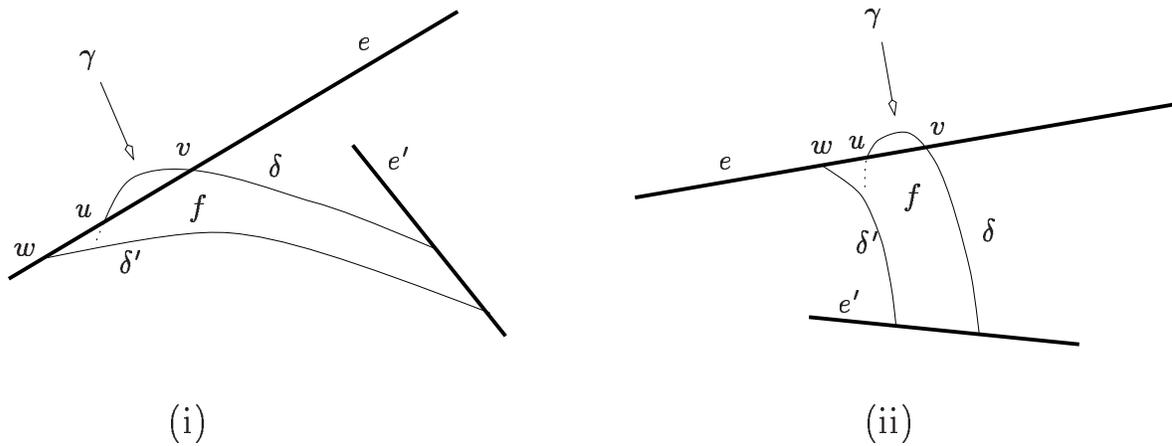


Figure 4: The portion of  $\gamma$  to the left of  $\gamma \cap e$  must lie fully within  $f$

(i)  $\delta$  and  $\delta'$  leave both  $e$  and  $e'$  from their bottom sides; see Figure 4(i). It is easily seen that in this case the left endpoints of  $\delta, \delta'$  lie on the same edge, say  $e$ . Suppose, without loss of generality, that  $v = \delta \cap e$  lies to the right of  $w = \delta' \cap e$ . Let  $u$  be the other intersection of  $\gamma$  with  $e$ ; by construction, such a point must exist. It is easily seen that  $u$  must lie between  $w$  and  $v$ , and that the extension of  $\gamma$  to the left of  $u$  lies in  $f$  locally near  $u$ . Moreover, this extension lies fully below the lines containing  $e$  and  $e'$ , respectively, implying that it cannot intersect any of the four sides of  $f$ , and thus the left endpoint of  $\gamma$  must lie inside  $f$ . The number of such faces is therefore at most  $m$ .

(ii)  $\delta$  and  $\delta'$  leave, say,  $e$  from its bottom side and  $e'$  from its top side; see Figure 4(ii). In this case the left endpoints of  $\delta$  and  $\delta'$  lie on  $e$ . Again, we assume that  $v = \delta \cap e$  lies to the right of  $w = \delta' \cap e$ . Let  $u$  be the other intersection of  $\gamma$  with  $e$ ; as above, by construction, such a point must exist,  $u$  must lie between  $w$  and  $v$ , and the extension of  $\gamma$  to the left of  $u$  lies in  $f$  locally near  $u$ . As above, we claim that this extension lies fully within  $f$ . As in the previous case, this extension cannot intersect  $e, \delta$ , or  $\delta'$ . Moreover, it also cannot meet  $e'$  (from its top side), since this would contradict our assumption that  $\delta$  is an edge of  $G$ . Hence, in this case too the left

endpoint of  $\gamma$  must lie inside  $f$ , so the overall number of faces of  $G$  of degree 2 is at most  $m$ .

Hence, a straightforward application of Euler's formula for planar maps implies that the number of edges of  $G$  is  $O(r^2 + m)$ .

Let  $m_\tau$  denote the number of short and bottom arcs in  $\tau$ . The preceding analysis implies that  $\sum_\tau m_\tau = O(r^2 + m)$ . The weaker bound obtained above implies that the number of mixed edge-touching regular vertices on  $\partial U$  that are formed within a trapezoid  $\tau$  by its  $m_\tau$  short and bottom arcs is  $O(m_\tau(n/r)^{1/2} + n/r)$ , for a total of

$$O\left(\sum_\tau \left[ m_\tau \left(\frac{n}{r}\right)^{1/2} + \frac{n}{r} \right]\right) = O\left((m + r^2) \left(\frac{n}{r}\right)^{1/2} + nr\right).$$

As argued above, the overall number of all other mixed edge-touching regular vertices on  $\partial U$  is  $O(m + nr)$ . If we choose  $r = \lceil m^{2/3}/n^{1/3} \rceil = \lceil n^{1/3} \rceil$ , the overall bound becomes  $O(n^{4/3})$ . This, combined with the lower bound construction in [15], completes the proof of Theorem 1.1.  $\square$

### 3 Transforming the Regions

The next lemma is used as a first step in the proofs of Theorems 1.2 and 1.3.

**Lemma 3.1** *Let  $\mathcal{C}$  be a collection of  $n$  Jordan regions in the plane, so that each pair of boundaries intersect in a finite number of points. Then we can transform this collection such that if  $c$  and  $c'$  are a pair of regions in  $\mathcal{C}$  that originally crossed regularly, where at least one of these two crossing points lies on  $\partial U$ , then after the transformation  $c$  and  $c'$  are openly disjoint and touch at a single point that lies on the new union boundary. Moreover, the number of intersections between any two region boundaries does not increase after the transformation. Finally, if all original regions in  $\mathcal{C}$  are convex, then they remain so after the transformation.*

**Proof:** The transformation process is iterative: Order the regions in  $\mathcal{C}$  arbitrarily as  $(c_1, \dots, c_n)$ . Let  $\mathcal{C}_i$  denote the collection after the first  $i$  steps of the transformation, with  $\mathcal{C}_0 = \mathcal{C}$ . To transform  $\mathcal{C}_{i-1}$  into  $\mathcal{C}_i$ , we do the following.

Put  $c = c_i$ , and apply a homeomorphism  $\tau_c$  of the plane that maps  $c$  onto the closed unit disk  $D$ ; the existence of such a homeomorphism is a consequence of Schönflies' Theorem [13]. If all the sets are convex, then we take  $\tau_c$  to be the identity and for uniformity put  $D = c$ . We now apply the following steps:

- (i) For each  $c' \neq c \in \mathcal{C}$ , we replace each maximal connected arc of  $D \cap \partial\tau_c(c')$  by the straight segment connecting its endpoints.
- (ii) We shortcut each maximal connected arc  $\gamma$  of  $\partial D \cap \text{int}(\tau_c(U))$  as follows. Let  $u$  and  $v$  denote the endpoints of  $\gamma$ , and let  $a, b \in \mathcal{C}$ ,  $a, b \neq c$ , be the sets whose transformed boundaries contain  $u, v$ , respectively. If  $a \neq b$ , we replace  $\gamma$  by the straight segment  $uv$ ; see the left side of Figure 5(iii). If  $a = b$ , we choose a point  $v'$  on  $\partial D$  that lies outside very close to  $v$  (and does not lie inside any other transformed set), and again replace the portion of  $\partial D$  between  $u$  and  $v'$  by the straight segment  $uv'$ ; see the right side of Figure 5(iii).
- (iii) We now transform the plane back using  $\tau_c^{-1}$ .

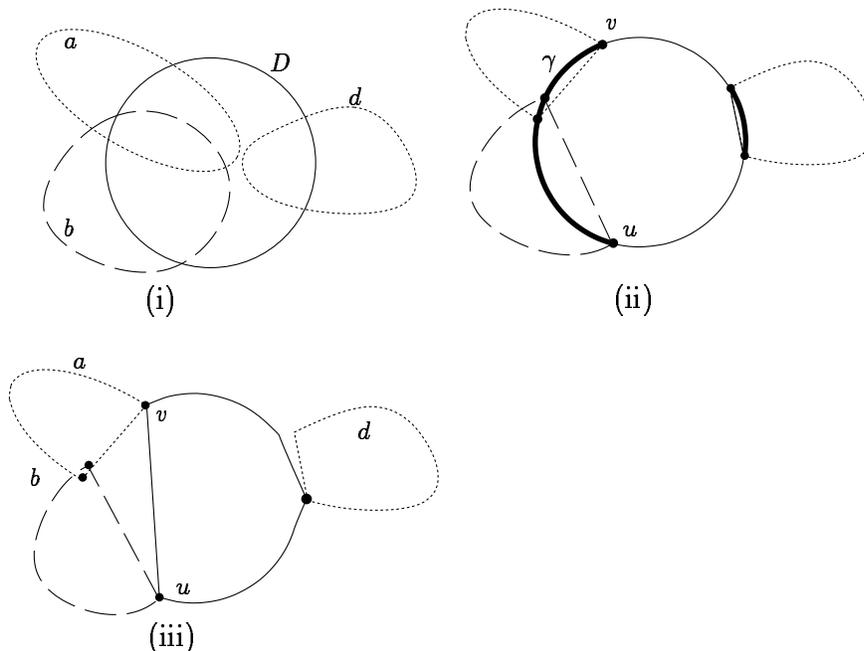


Figure 5: Demonstration of the transformation rules: (i) The disk  $D$  with the images of three other sets, referred to as  $a, b$ , and  $d$ . (ii) The modified sets  $a, b$  and  $d$ ; the arcs  $\gamma \subseteq \partial D$  that are shortcut in the next stage, are highlighted. (iii)  $D$  after applying stage (ii) of the transformation.

We iterate this procedure, applying it to each  $c_i \in \mathcal{C}$  in the above order.

We claim that this transformation does not increase the number of intersections between any pair of boundaries. Indeed, consider the step where a set  $c_i$  is processed, and let  $a, b$  be two distinct sets in  $\mathcal{C}_{i-1}$ . If  $c_i \neq a, b$  then the portions of  $\partial a$  and  $\partial b$  outside  $c$  do not change, while their portions inside  $c_i$ , after applying  $\tau_{c_i}$  and our “straightening” step, consist of straight segments. If two such segments  $uv$  and  $wz$  lying on the modified respective boundaries  $\partial a, \partial b$  inside  $D$ , cross at some point  $x$ , then, since  $u, w, v, z$  all lie on  $\partial D$  and must appear there in this cyclic order, and since the original respective portions of  $\partial \tau_{c_i}(a)$  and  $\partial \tau_{c_i}(b)$  connecting  $u$  to  $v$  and  $w$  to  $z$  lie fully inside  $D$ , it is easy to verify that these original portions must cross each other at some point  $x'$ . We charge  $x$  to  $x'$ , and note that this charging is unique, implying that the number of intersections has not increased.

Suppose then that  $c_i = a$ , say, and let  $x$  be a new intersection between  $\partial D$  and  $\partial \tau_{c_i}(b)$ , after the straightening step. Then  $x$  must lie (in the  $\tau_a$  image) on one of the new straight segments  $uv$  or  $uv'$  forming  $\partial a$  in the transformed plane, and on some new straight portion  $wz$  of the transformed  $\partial b$ . If  $x$  is a common endpoint of these two straight segments, then it must have been an intersection point between  $\partial a$  and  $\partial b$  in  $\mathcal{C}_{i-1}$ , so no new intersection arises in this case. Otherwise, as above, it follows that one of  $w, z$ , say  $w$ , must lie on the portion of  $\partial D$  that has been replaced by  $uv$  or  $uv'$ , so we can charge  $x$  to  $w$ , which was an old intersection point of  $\partial a$  and  $\partial b$  (in the transformed plane), and this charging is unique, again implying that the number of intersections has not increased.

In particular, every pair of new boundaries intersect (transversally) at most  $s$  times, and every pair of boundaries that originally intersected in two (regular) vertices either continue to do so, or just touch each other in a single point, or do not intersect at all after the transformation. Specifically, in the case of regular intersection between two boundaries  $\partial a, \partial b$ , if at least one of the two intersections lies on  $\partial U$ , then the two transformed boundaries touch at a single point that lies on  $\partial U$ ; see Figure 5(iii). Indeed, suppose that  $\partial a$  and  $\partial b$  intersect at exactly two points  $u, v$ , one of which lies on  $\partial U$ , and suppose that  $a$  is processed before  $b$ . When  $a$  is processed,  $bda$  and  $\partial b$  become touching by construction, and it is easy to verify that this situation does not change after any subsequent transformation step (one needs to verify this only for the step that processes  $b$  because no other step affects the neighborhood of this touching point).

All these considerations complete the proof of the lemma.  $\square$ .

Observe that  $U$  might be changed by these transformations: The original  $\partial U$  is still a portion of the new boundary, but  $U$  could have gained additional ‘holes’, as shown in Figure 5(iii). Note also that the result of the transformation may depend on the order in which the sets are processed.

## 4 Convex Regions: Proof of Theorem 1.2

We first apply to  $\mathcal{C}$  the transformation described in the previous section. As already noted, if the sets are convex, then there is no need to apply a homeomorphism to the plane at each step of this process, and the sets remain convex after the transformation. Recall that this transformation has the property that if  $c$  and  $c'$  are a pair of sets in  $\mathcal{C}$  that originally crossed regularly, where at least one of these two crossing points lies on  $\partial U$ , then after the transformation  $c$  and  $c'$  are openly disjoint and touch at a single point that lies on the new union boundary. (If  $c$  and  $c'$  met regularly but both intersection points lay in the interior of  $U$ , then the transformed sets become disjoint.) The converse is also true: Any boundary touching between two transformed sets must occur on  $\partial U$ . Also, the new sets retain the property that any pair of boundaries cross in at most  $s$  points. From now on, we assume that  $\mathcal{C}$  has the properties just noted.

For each  $c \in \mathcal{C}$ , let  $e_c$  denote its ‘equator’, namely the segment connecting the leftmost point and the rightmost point in  $c$ . (By an appropriate general position assumption, or by appropriate tilting of the plane,  $e_c$  is uniquely defined for each  $c \in \mathcal{C}$ ; these segments were called ‘sentinels’ in [1].)

We first construct a *hereditary* segment tree  $Q$  on the  $x$ -projections of the equators, as in [3] (consult [3] for more details, and for the terminology that we use below). Each node  $v$  of  $Q$  stores the standard segment-tree list  $L_v$  of so-called ‘long’ equators, and also a list  $S_v$  of ‘short’ equators, those that are stored in some list  $L_w$ , for a proper descendant  $w$  of  $v$ , and thus have at least one endpoint in the interior of the vertical strip  $\sigma_v$  associated with  $v$ . It is easily verified that for any boundary touching between two sets  $a, b \in \mathcal{C}$ , there is a unique node  $v$  of  $Q$  (on the path to the leaf  $w$  whose strip  $\sigma_w$  contains the touching), such that the touching point lies in  $\sigma_v$  and either both  $a$  and  $b$  belong to  $L_v$  or one of them belongs to  $L_v$  and the other to  $S_v$ . We have  $\sum_v (|L_v| + |S_v|) = O(n \log n)$ . Also, any such touching occurs between the upper boundary of one set and the lower boundary of the other.

We now fix a node  $v$  and bound the number of boundary touchings within  $\sigma_v$  formed between two sets in  $L_v \cup S_v$ , at least one of which lies in  $L_v$ . We will only describe the case where the

other set lies in  $S_v$ , because the case where both of them lie in  $L_v$  is simpler and can be handled by a similar approach. Moreover, with no loss of generality, it suffices to consider only boundary touchings where the set in  $L_v$  lies above the set in  $S_v$ .

We apply fairly standard range-searching techniques to obtain a finite collection  $\{A_i \times B_i\}_i$  of complete bipartite graphs, such that

- (a) For each  $i$ ,  $A_i \subseteq S_v$  and  $B_i \subseteq L_v$ .
- (b) For each  $i$  there is a substrip  $\sigma^{(i)} \subseteq \sigma_v$ , such that for each  $a \in A_i$ ,  $b \in B_i$ , the  $x$ -projection of  $a$  contains that of  $\sigma^{(i)}$ , the equator  $e_b$  lies fully above  $a$  within  $\sigma_v$ , and the equator  $e_a$  lies fully below  $b$  within  $\sigma^{(i)}$ . See Figure 6.
- (c) For each pair of sets  $a \in S_v$ ,  $b \in L_v$ , such that  $\partial a$  touches  $\partial b$  at a point  $\zeta$  within  $\sigma_v$  and  $a$  lies below  $b$ , there is an index  $i$  such that  $a \in A_i$ ,  $b \in B_i$ , and  $\zeta \in \sigma^{(i)}$ .
- (d)  $\sum_i (|A_i| + |B_i|) = O(|S_v||L_v|^{1/2+\epsilon} + |L_v||S_v|^{1/2+\epsilon})$ , for any  $\epsilon > 0$ , where the constant of proportionality depends on  $\epsilon$  and on  $s$ .

Suppose this has been done. Then fix an index  $i$ , and note that any boundary touching  $\zeta$  that occurs within  $\sigma^{(i)}$  between a set  $a$  in  $A_i$  and a set  $b$  in  $B_i$  must occur along the upper envelope of the upper boundaries of the sets of  $A_i$ , and along the lower envelope of the lower boundaries of the sets of  $B_i$ . Indeed, suppose to the contrary that, say,  $\zeta$  does not lie on this upper envelope. Since  $\zeta$  lies on the boundary of the union of  $A_i$ , the set  $a'$  appearing on the envelope at the  $x$ -coordinate of  $\zeta$  must be such that  $e_{a'}$  lies above  $\zeta$ , which contradicts the property that  $e_{a'}$  has to lie fully below  $b$  within  $\sigma^{(i)}$ . The argument for the lower envelope is fully symmetric. The number of such touchings is therefore proportional to the combined complexity of these envelopes within  $\sigma^{(i)}$ , which is at most  $\lambda_s(|A_i|) + \lambda_s(|B_i|)$ , where  $\lambda_s(n)$  denotes the maximum length of  $(n, s)$ -Davenport-Schinzel sequences [16]. Summing over all  $i$ 's, we obtain that the total number of boundary touchings 'associated' with  $v$  is proportional to  $O(|S_v||L_v|^{1/2+\epsilon} + |L_v||S_v|^{1/2+\epsilon})$ , for a slightly larger, but still arbitrarily small  $\epsilon > 0$ . Summing these bounds over all nodes  $v$  of  $Q$ , we obtain the overall bound  $O(n^{3/2+\epsilon})$ , again for slightly larger, but still arbitrarily small  $\epsilon > 0$ , with the constant of proportionality depending on  $\epsilon$  and on  $s$ . This therefore completes the proof of the theorem.

To obtain the decomposition  $\{A_i \times B_i\}$ , we use a multi-level range-searching structure, where each node  $w$  at each level of the structure will store a complete bipartite graph  $A'_w \times B'_w$ , such that  $A'_w \subseteq S_v$  and  $B'_w \subseteq L_v$ . Each subsequent level of the structure 'enforces' some more of the desired constraints, and the bipartite graphs within each subsequent level form a refinement of the graphs obtained at the previous level. See [2] for more details concerning multi-level range-searching structures.

In the first level we enforce the property that the  $e_b$ 's, for  $b \in L_v$ , lie above the  $a$ 's in  $S_v$ . In what follows, we clip all the relevant  $a$ 's and  $b$ 's to within  $\sigma_v$ . Actually, we want to have the property that the equator  $e_b$  of any such  $b$  lie above the upper boundary of any such  $a$ .

We may replace any  $e_b$  by the line containing it (this has no effect on what happens within  $\sigma_v$ ), and replace any  $a$  by the portion of  $\sigma_v$  lying below the upper boundary of the original  $a$  (so we first make  $a$  smaller by ignoring its portion outside  $\sigma_v$ , and then make  $a$  larger by allowing it to expand downwards within  $\sigma_v$ ). We apply a standard duality to the plane (as in [5]), which preserves incidences and the above/below relationship. This duality maps the upper boundary of any  $a \in S_v$  to a convex  $x$ -monotone curve  $\gamma_a$ , and the (extended) equators  $e_b$  of sets in  $L_v$  are mapped to points  $e_b^*$ . An equator  $e_b$  lies above the upper boundary of  $a$  if and only if the dual point  $e_b^*$  lies above  $\gamma_a$ . Note also that any pair of curves  $\gamma_a, \gamma_{a'}$  intersect each other at most  $s$  times, because any such intersection point is the dual of a common tangent to the upper boundaries of  $a$  and  $a'$ , and there can be at most  $s$  such common tangents, because  $\partial a$  and  $\partial a'$  intersect in at most  $s$  points.

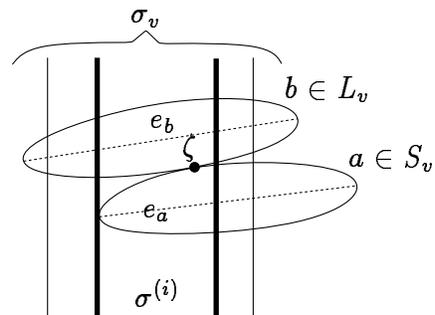


Figure 6: An illustration of property (b).

Thus, in this dual setting, the desired first-level decomposition  $\{A'_w \times B'_w\}_w$  of  $S_v \times L_v$  has to satisfy the following properties:

- (e) For each  $w$ , the point  $e_b^*$  dual to the equator  $e_b$  of any set  $b \in B'_w$  lies above  $\gamma_a$ , for every  $a \in A'_w$ .
- (f) For any  $a \in S_v, b \in L_v$ , such that  $e_b^*$  lies above  $\gamma_a$ , there is a  $w$  such that  $a \in A'_w$  and  $b \in B'_w$ .

Put  $m_v = |S_v|$ ,  $n_v = |L_v|$ . To achieve this decomposition, we fix some sufficiently large constant parameter  $\xi$ , draw a random sample  $R$  of  $\xi$  sets  $a \in S_v$ , consider the arrangement  $\mathcal{A}_R$  of the corresponding curves  $\gamma_a$ , and apply a vertical decomposition to  $\mathcal{A}_R$  that produces  $O(\xi^2)$  pseudo-trapezoidal cells. Since each pseudo-trapezoid is determined by at most four curves  $\gamma_a$ , it follows from [4] that with high probability, no pseudo-trapezoid is crossed by more than  $\frac{cm_v}{\xi} \log \xi$  curves, for some appropriate constant  $c$ . We may assume that our sample  $R$  does indeed have this property.

For each pseudo-trapezoid  $\tau$ , let  $B_\tau$  denote the subset of  $L_v$  consisting of those sets  $b$  whose dual points  $e_b^*$  lie inside  $\tau$ , and let  $A_\tau$  (resp.  $C_\tau$ ) denote the subset of  $S_v$  consisting of those sets  $a$  whose dual curves  $\gamma_a$  pass below  $\tau$  (resp. cross  $\tau$ ). Put  $k_\tau = |A_\tau|$ ,  $n_\tau = |B_\tau|$ , and  $m_\tau = |C_\tau|$ . We have  $\sum_\tau n_\tau = n_v$ ,  $k_\tau \leq m_v$ , and  $n_\tau \leq \frac{cm_v}{\xi} \log \xi$ . By partitioning  $\tau$  vertically into subcells, if necessary, we may also assume that  $n_\tau \leq n_v/\xi^2$  for each  $\tau$ , while the total number of subcells remains  $O(\xi^2)$ .

We add to the first-level output collection of complete bipartite graphs all the products  $A_\tau \times B_\tau$ , and repeat the whole process recursively within each cell  $\tau$ , with the sets  $C_\tau$  and  $B_\tau$ . We stop the recursion when the size of  $C_\tau$  or of  $B_\tau$  falls below some specified constant, and then output all appropriate singleton products  $\{a\} \times \{b\}$ , for  $a \in C_\tau, b \in B_\tau$ .

It is clear from the construction that the resulting decomposition satisfies the required properties (e) and (f). We next estimate its total size  $\sum_\tau (|A_\tau| + |B_\tau|)$ . At the top level of the recursion we have  $\sum_\tau |B_\tau| = n_v$ . Since the number of levels of recursion is  $O(\log n_v)$  and the points  $e_b^*$  are partitioned among the recursive subproblems, it follows that in total we have

$\sum_{\tau} |B_{\tau}| = O(n_v \log n_v)$ . Similarly, at the top level of the recursion we have  $\sum_{\tau} |A_{\tau}| = O(m_v \xi^2)$ , and  $\sum_{\tau} |C_{\tau}| = O(m_v \xi \log \xi)$  (where both constants of proportionality depend only on  $s$ ). The maximum depth  $j$  of the recursion satisfies  $\xi^{2j} \leq n_v$ , or  $\xi^j \leq n_v^{1/2}$ . It follows that, for an appropriate constant  $c'$  that depends on  $s$ , the overall sum  $\sum_{\tau} |A_{\tau}|$  is at most proportional to

$$m_v \xi^2 \left(1 + c' \xi \log \xi + \cdots + (c' \xi \log \xi)^{j-1}\right) = O(m_v (c' \xi \log \xi)^j) = O(m_v n_v^{1/2+\varepsilon}),$$

where we can make  $\varepsilon > 0$  arbitrarily small by choosing  $\xi$  sufficiently large (as a function of  $\varepsilon$  and  $s$ ). Hence we have

$$\sum_{\tau} |A_{\tau}| = O(m_v n_v^{1/2+\varepsilon}) \quad \text{and} \quad \sum_{\tau} |B_{\tau}| = O(n_v \log n_v). \quad (1)$$

Next, fix a pair  $A = A_{\tau}$ ,  $B = B_{\tau}$  in this decomposition. For each  $a \in A$ ,  $b \in B$  the equator  $e_b$  lies above the upper boundary of  $a$  (within  $\sigma_v$ ). This already implies (arguing as above) that any boundary touching  $\zeta$  between  $\partial a$  and  $\partial b$ , for any pair  $a \in A$ ,  $b \in B$  that lies on the union, must lie on the lower envelope  $E$  of the lower boundaries of the sets in  $B$ . The complexity of  $E$  is at most  $\lambda_s(|B|)$ , which is nearly linear in  $|B|$ .

For each  $a \in A$ , let  $a_x$  denote its  $x$ -projection, clipped to within  $\sigma_v$ . We construct a secondary segment tree  $R$  on the intervals  $a_x$ . Each node  $u$  of  $R$  is associated with a vertical strip  $\sigma^{(u)} \subseteq \sigma_v$ . In addition to the standard segment-tree list  $A^{(u)}$  of  $x$ -projections of sets in  $A$  that is stored at  $u$ , we also store there a list  $B^{(u)}$  of the sets in  $B$  whose lower boundaries appear in  $E \cap \sigma^{(u)}$ . We clearly have

$$\sum_u |A^{(u)}| = O(|A| \log |A|)$$

and

$$\sum_{|B^{(u)}| > 1} |B^{(u)}| = O(\lambda_s(|B|) \log |A|).$$

Indeed, the first bound is a standard property of segment trees. The second bound is a consequence of the following observations: (a) The number of breakpoints of  $E$  is at most  $\lambda_s(|B|)$ . (b) Each breakpoint belongs to at most  $\log |A|$  strips  $\sigma^{(u)}$ . (c) For any  $u$ , the size of  $B^{(u)}$  is upper bounded by 1 plus the number of breakpoints of  $E$  in  $\sigma^{(u)}$ .

Now suppose that there is a pair  $a \in A$ ,  $b \in B$ , with a touching  $\zeta$  between the upper part of  $\partial a$  and the lower part of  $\partial b$  that lies on the union. There is a unique node  $u$  of  $R$  such that  $a \in A^{(u)}$  and  $\zeta \in \sigma^{(u)}$ . Hence  $\partial b$  appears on  $E$  within  $\sigma^{(u)}$ , so  $b \in B^{(u)}$ . Note also that the line containing  $e_a$  passes fully below  $b \cap \sigma^{(u)}$  (because the  $x$ -projection of  $a$  fully contains the projection of  $\sigma^{(u)}$ ).

We now fix a node  $u$  for which  $|B^{(u)}| > 1$ , and apply a symmetric version of the first-level decomposition to  $A^{(u)} \times B^{(u)}$ , to obtain a collection  $\{A'_j \times B'_j\}_j$  of complete bipartite graphs, such that

- (g) For each  $j$ , we have  $A'_j \subseteq A^{(u)}$  and  $B'_j \subseteq B^{(u)}$ ; we also associate the strip  $\sigma^{(u)}$  with  $j$ .
- (h) For each  $j$ , each  $a \in A'_j$  and each  $b \in B'_j$ , the equator  $e_a$  lies fully below  $b$  within  $\sigma^{(u)}$ .
- (i) For any boundary touching  $\zeta$  as above, there is an index  $j$  such that  $a \in A'_j$  and  $b \in B'_j$  (and  $\zeta$  lies in the strip of  $j$ , that is, in  $\sigma^{(u)}$ ).

$$(j) \sum_j (|A'_j| + |B'_j|) = O(|A^{(u)}| \log |A^{(u)}| + |B^{(u)}| |A^{(u)}|^{1/2+\varepsilon}), \text{ for any } \varepsilon > 0.$$

If  $|B^{(u)}| = 1$  we output only one bipartite graph  $\tilde{A}^{(u)} \times B^{(u)}$  where  $\tilde{A}^{(u)}$  consists of all  $a \in A^{(u)}$  such that  $e_a$  lies fully below  $b$  within  $\sigma^{(u)}$ .

The grand collection of complete bipartite graphs  $A'_j \times B'_j$ , gathered over all nodes  $u$  of  $R$ , and over all first-level pairs  $A_\tau \times B_\tau$ , is the desired output collection. It clearly satisfies properties (a)–(c). Concerning (d), we first sum the bounds (j) over all  $u \in R$ , to obtain the bound

$$O(|A| \log^2 |A| + |B| |A|^{1/2+\varepsilon}),$$

for a slightly larger but still arbitrarily small  $\varepsilon > 0$ , where the constant of proportionality depends on  $\varepsilon$  and on  $s$ . Clearly, this bound also subsumes the case  $|B^{(u)}| = 1$ .

Finally, we sum these bounds over all pairs  $A = A_\tau, B = B_\tau$ , in the first-level decomposition, and make use of (1), to conclude that this sum is at most proportional to

$$\sum_\tau O(|A_\tau| \log^2 |A_\tau| + |B_\tau| |A_\tau|^{1/2+\varepsilon}) = O(m_v n_v^{1/2+\varepsilon} + n_v m_v^{1/2+\varepsilon}),$$

again, for a slightly larger but still arbitrarily small  $\varepsilon > 0$ , where the constant of proportionality depends on  $\varepsilon$  and on  $s$ . Hence (d) is also satisfied. As already noted, this completes the proof of the theorem.  $\square$

**Remark:** An obvious open problem is to close the gap between this upper bound and the lower bound  $\Omega(n^{4/3})$  noted earlier.

## 5 Arbitrary Regions: Proof of Theorem 1.3

We prove the theorem using the following ‘forbidden subgraph’ argument. We first transform  $\mathcal{C}$  as described in Lemma 3.1, but continue to denote the transformed collection by  $\mathcal{C}$ . Now we define a graph  $H$  whose vertices are the regions in  $\mathcal{C}$  and whose edges connect pairs of regions that touch at a point on the boundary of the union. We claim that  $H$  does not contain a complete bipartite graph  $K_{r,g}$  for some sufficiently large constant  $r = r(s), g = g(s)$  that depend on  $s$ , where  $s$  is the maximum number of intersections between any two region boundaries.

Indeed suppose to the contrary that there exist subsets  $\mathcal{R}, \mathcal{G}$  of  $\mathcal{C}$  of sizes  $r, g$ , respectively, such that for each  $(c, c') \in \mathcal{R} \times \mathcal{G}$ , the boundaries of  $c$  and  $c'$  touch at a point that lies on the boundary of the union  $U$ . We claim that this is impossible for sufficiently large values of  $r$  and  $g$  that depend only on  $s$ . With no loss of generality, we may assume that  $\mathcal{C} = \mathcal{R} \cup \mathcal{G}$ . We refer to sets of  $\mathcal{R}$  (resp.  $\mathcal{G}$ ) as “red” (resp. “green”).

Consider the arrangement  $\mathcal{A}(\mathcal{G})$  and let  $U_{\mathcal{G}}$  denote the union of  $\mathcal{G}$ . We claim that any  $c \in \mathcal{R}$  is fully contained in the closure of a single hole (connected component of the complement) of  $U_{\mathcal{G}}$ . Indeed, if this were false,  $\partial c$  would have to properly cross some boundary of a green region  $c'$ , which is impossible by construction.

Consider the collection of holes of  $U_{\mathcal{G}}$  that contain red regions. Since each red region must touch the boundary of every green set at a point that lies on  $\partial U_{\mathcal{G}}$ , it follows that all these holes are part of the zone in  $\mathcal{A}(\mathcal{G})$  of any green boundary. Hence the overall complexity of all these holes is at most  $a\lambda_{s+2}(g)$ , for some absolute constant  $a$  (see [14]). Similarly, by interchanging

$\mathcal{G}$  and  $\mathcal{R}$ , the overall complexity of the holes of the red union  $U_{\mathcal{R}}$  that contain green regions is at most  $a\lambda_{s+2}(r)$ .

We now construct a planar bipartite graph  $G$ , whose nodes are the arcs forming the boundaries of these green holes of  $U_{\mathcal{G}}$  and red holes of  $U_{\mathcal{R}}$ ; each edge of  $G$  connects a green arc  $\zeta$  to a red arc  $\eta$  if these two arcs touch at a regular vertex of the overall union. It is easy to verify that  $G$  is indeed planar, and that it has no multiple edges. Since the graph is bipartite, the number of its edges is at most twice the number of its nodes, that is, at most  $2a(\lambda_{s+2}(g) + \lambda_{s+2}(r))$ . On the other hand, the number of edges of  $G$  must be  $rg$ , because every green set and every red set touch at some regular vertex on  $\partial U$ . We thus obtain:

$$rg \leq 2a(\lambda_{s+2}(g) + \lambda_{s+2}(r)),$$

which is impossible if  $r$  and  $g$  are sufficiently large constants (that depend on  $s$ ).

Using standard results in extremal graph theory (see [14]), it follows that the number of regular vertices on the boundary of the union of  $\mathcal{C}$  is  $O(n^{2-1/r})$ . This completes the proof of the theorem.  $\square$

**Remarks.** (1) The following is a natural extension of the concept of regular intersection: We say that the boundaries of two Jordan regions  $a$  and  $b$  meet regularly at  $u$  and  $v$  if there is a connected component of  $a \cap b$  that contains  $u$  and  $v$  and no other intersection point of  $\partial a \cap \partial b$  on its boundary. In other words, we allow  $\partial a$  and  $\partial b$  to cross more than twice, but require that their intersections  $u$  and  $v$  behave ‘locally’ as regular intersections; see Figure 7. Unfortunately, it is possible to construct families of  $n$  Jordan regions, each pair of whose boundaries cross at most 6 times, such that the boundary of their union has  $\Theta(n^2)$  regular vertices under this extended definition. Such a construction is depicted in Figure 8 where  $n$  sets creating  $\Theta(n^2)$  regular vertices (under the new definition) on their union. We start with  $n/2$  long and thin ‘horizontal’ rectangles and  $n/2$  long and thin ‘vertical’ rectangles that form a grid. For each of the  $\Theta(n^2)$  bounded holes of their union, take the rectangle  $R$  whose top side forms the bottom edge of the hole, and deform this side so as to create a small protrusion that extends into the hole, bends to the right, and meet the right side of the hole at two points, which are regular according to the new definition. Clearly, any pair of boundaries now cross at most 6 times. Thus Theorem 1.3 fails for this extended notion of regularity.

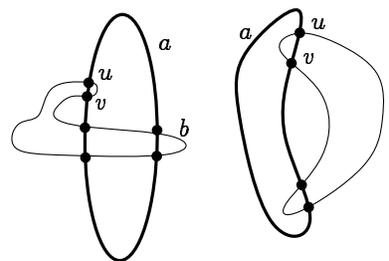


Figure 7: According to the extended notion of regularity, both pairs  $u, v$  are regular

In the lower bound construction just given, even though each pair of boundaries cross at most 6 times, the shape of some of the boundaries is quite complicated. We do not know whether a similar construction can be obtained for semialgebraic Jordan regions of constant description complexity<sup>2</sup>, so the question whether Theorem 1.3 holds in this case remains open.

(2) Nevertheless, if any two boundaries cross at most 4 times, then if two boundaries cross four times but meet regularly, according to the new definition, then the four intersections come in two pairs, each of which is a regular pair; see Figure 7(right). It is easy to verify that Theorem 1.3 continues to hold in this case, using the same proof with slight and obvious modifications.

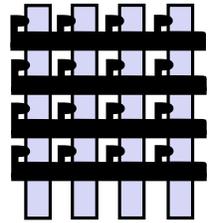


Figure 8:

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<sup>2</sup>That is, regions that are each described by a boolean combination of a constant number of polynomial equalities and inequalities of constant maximum degree.

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